



# Image decompositions and transformations as peaks and wells

Fernand Meyer

## ► To cite this version:

Fernand Meyer. Image decompositions and transformations as peaks and wells. 10th International Symposium on Mathematical Morphology and Its Application to Signal and Image Processing, ISMM 2011, Jul 2011, Verbania-Intra, Italy. pp.25-36, 10.1007/978-3-642-21569-8\_3 . hal-00547277

**HAL Id: hal-00547277**

**<https://hal.science/hal-00547277>**

Submitted on 15 Dec 2010

**HAL** is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

# Image decompositions and transformations as peaks and wells

Fernand Meyer

Mines ParisTech,  
Département maths et systèmes,  
Centre de Morphologie Mathématique,  
F-77305 Fontainebleau Cedex, France

**Abstract.** An image may be decomposed as a difference between an image of peaks and an image of wells. Applying a morphological operator to these two components before reconstructing a final image produces interesting filters for grey tone or binary images. This decomposition depends upon the point of view from where the image is considered.

## 1 Introduction

A binary image is made of particles and holes. Each particle may contain one or several holes and each hole one or several particles. These structures may be deeply nested. For describing this structure, J.Serra [5] introduced the homotopy tree, H.Heijmans [2] called it the adjacency tree. R. Keshet [3] and C. Ballester [1] studied it in depth and gave algorithms for constructing it. They presented also each a different extension to grey-tone images of this tree, resulting in the so called tree of shapes.

Here, we propose a description of images which is not as complete as the tree of shapes, but which is extremely simple and fast to construct. Given a reference set  $X$ , called view-point set, we decompose an image in a difference between two components, one representing its peaks, the other its wells. They represent respectively the sums of positive and negative variations of the image if one follows a paths starting in  $X$ . Both components belong to the lattice of images without minima outside  $X$ , on which we define an adjunction.

Ch.Ronse proposed an identical decomposition for functions defined strictly increasing elements of a poset  $P$ , with the final aim to find a sound way for constructing a flat operator on gray level image from a non-increasing operator on binary images [4].

In the present paper we decompose functions along connected paths in graphs. Applying a morphological operator to these two components before reconstructing a final image produces interesting filters for grey tone or binary images. This decomposition depends upon the point of view from where the image is considered.

## 2 Scheduling trains

### 2.1 Setting the scenario of trains circulating on graphs

Let  $G = (N, E)$  be a graph where nodes  $N$  and oriented edges or arrows  $E$  are weighted with positive weights. The graph may be considered as a railway network, where the nodes are railway stations and the arrows are connections between them. Trains may follow all possible paths on  $G$ . However, they may only start from a subset  $X \subset N$  of railway stations.

Consider a particular train. It starts at station  $s \in X$  at time  $\tau(s)$ , follows a path  $\theta = (x_0 = s, x_1, \dots, x_n = t)$  where  $x_i$  and  $x_{i+1}$  are two railway stations linked by an arrow  $(i, i+1)$  of  $E$ , weighted by the time  $e_{i,i+1}$  needed for following it. The arrival time at destination is then  $\tau(s) + \sum_{i,i+1 \in \theta} e_{i,i+1}$ .

We now consider all trains starting at all possible nodes and taking all possible routes, and observe the earliest time when each node is reached by a train; if a train arrives at  $i$  before  $\tau(i)$ , we replace  $\tau(i)$  by this first arrival time. For some nodes no train ever arrives;  $i$  is such a node, if  $\tau(i)$  is too early, so that all trains coming from another node arrive at  $i$  after  $\tau(i)$ . For others, no train departs: if  $\tau(i)$  is too late, no train starting from  $i$  has a chance to be the first to reach another node; this is in particular the case if  $\tau(i) = \infty$ . Some nodes cumulate both situations, and no train departs or arrives.

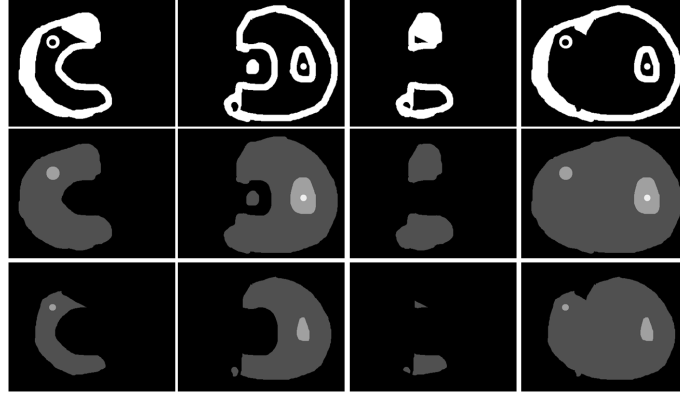
The resulting schedules  $\hat{\tau} = \Theta_X(e, \tau)$  depends on the distribution of initial departure times and crossing times of each edge. Defining  $\tau_X^\infty$  by  $\tau_X^\infty = \tau$  on  $X$  and  $\tau_X^\infty = \infty$  elsewhere, it is obvious that  $\Theta_X(e, \tau) = \Theta_N(e, \tau_X^\infty)$ , since no train with an infinite departure time has the chance to reach another node first.  $\Theta_N(e, \tau_X^\infty)$  is clearly an opening on the initial distribution of departure times on all nodes; it is obviously anti-extensive and increasing. It is also idempotent, as a second scheduling would not change the distribution  $\tau(i)$  any further. We also remark that if  $\tau(s) = 0$  for  $s \in X$ , the resulting schedules simply are the shortest path between  $X$  and all other nodes.

### 2.2 Harmonizing the schedules is a shortest path problem in a completed graph

**Completed graph and extended functions** We now fix  $G, e$  and  $X$  and consider varying distributions of starting times. Any such distribution is a function defined on the nodes of  $G$ . In case of images defined on a grid,  $G$  is the graph associated to the grid: pixels are nodes and neighboring pixels are connected by an edge.

We obtain an augmented graph  $G_X$ , by adding to  $G$  a dummy node  $\Omega$  with weight  $\tau(\Omega) = 0$  and dummy edges  $(\Omega, i)$  between  $\Omega$  and each node  $i$  of  $X$ , with  $e_{\Omega i} = \tau_i$ . Since the travelling time along the edge  $(\Omega, i)$  is  $\tau_i$ , it is equivalent for a train to start from  $i$  at time  $\tau_i$  or from  $\Omega$  at time 0 and follow the edge  $(\Omega, i)$  for reaching  $i$ . The earliest time for a train to arrive at any node  $k$  of  $G$  is the total duration of the shortest path between  $\Omega$  and  $k$ . It follows that scheduling

the graph  $G$  amounts to constructing the shortest path between  $\Omega$  and all other nodes in the graph  $G_X$ , which is a classical problem in graph theory for which many algorithms exist. In "Scheduling trains with delayed departures" in the same issue, we have presented the algorithms of Moore Dijkstra and of Berge and shown that, if the introduction of a dummy node  $\Omega$  and dummy edges  $(\Omega, i)$  is a useful support for thinking, it is not necessary in practice.



**Fig. 1.**

Columns 1 and 2 : Binary sets  $X$  and  $Y$  and their decomposition

Column 3 :  $X \wedge Y$  and its decomposition

Column 4 :  $X \vee Y$  and its decomposition

### 3 Decomposing an image into peaks and wells

We consider now an image  $f$  defined on the the nodes  $N$  of the graph  $G$ . We define three distributions of edge weights on  $G$  and its extension  $G_X$ , derived from the distribution of nodes. We consider two nodes  $i$  and  $j$  of  $G$  and a node  $k \in X$ . The weights of  $e_{ij}$  in  $G$  and of  $e_{\Omega k}$  in  $G_X$  are:

- $e_{ij}^+ = f_{ij}^+ = \vee(f_j - f_i, 0)$  and  $e_{\Omega k}^+ = f_k$ , positive for upwards transitions and null otherwise.
- $e_{ij}^- = f_{ij}^- = \vee(f_i - f_j, 0)$  and  $e_{\Omega k}^- = 0$  positive for downwards transitions and null otherwise.
- $e_{ij}^\# = |f_j - f_i| = |f_{ij}|$  and  $e_{\Omega k}^\# = f_k > 0$ .

Consider now an arbitrary node  $x$  of  $G$  and an arbitrary path  $\pi = (\Omega, x_1, x_2, \dots, x_n)$  between  $\Omega$  and  $x$ . We may decompose  $f(x)$  along the path  $\pi$  as follows:

$$f(x) = f(x) - f(x_{n-1}) + f(x_{n-1}) - f(x_{n-2}) + \dots + f(x_2) - f(x_1) + f(x_1) - f(\Omega) + f(\Omega)$$

Since  $f_j - f_i = f_{ij}^+ - f_{ij}^-$  for neighboring pixels  $i, j$  in the path  $\pi$ , and  $f(\Omega) = 0$ , we get  $f(x) = f(x_1) + \sum_{ij \in \pi} f_{ij}^+ - \sum_{ij \in \pi} f_{ij}^-$ , where  $x_1 \in X$  is linked to  $\Omega$ .

Among all possible paths, there is a path  $\hat{\pi}$  for which  $\sum_{ij \in \pi} f_{ij}^-$  is minimal. As  $f(x_1) + \sum_{ij \in \pi} f_{ij}^+ - \sum_{ij \in \pi} f_{ij}^-$  has a constant value, the expression  $f(x_1) + \sum_{ij \in \pi} f_{ij}^+$  is also minimal on  $\hat{\pi}$  and so is their sum  $f(x_1) + \sum_{ij \in \pi} f_{ij}^+ + \sum_{ij \in \pi} f_{ij}^- = f(x_1) + \sum_{ij \in \pi} |f_{ij}|$ .  $f(x_1) + \sum_{ij \in \pi} f_{ij}^+$ ,  $\sum_{ij \in \pi} f_{ij}^-$  and  $f(x_1) + \sum_{ij \in \pi} |f_{ij}|$  respectively represent the positive, negative and total variation along  $\hat{\pi}$  in  $G_X$ .

The length  $f(x_1) + \sum_{ij \in \pi} f_{ij}^+$  of the shortest path between  $\Omega$  and  $x$  in the graph  $G$  with edge weights  $e^+$  is the schedule of the first train arriving at  $x$  in  $G_X$ . We write  $\Theta_X(e^+, f)$  for the distribution of schedules on all nodes. Likewise, the schedules for the distribution of weights  $e^+$  and  $e^\#$  are  $f_X^-(x) = \sum_{ij \in \pi} f_{ij}^- = \Theta_X(e^-, f)$  and  $f_X^\#(x) = f(x_1) + \sum_{ij \in \pi} f_{ij}^+ + \sum_{ij \in \pi} f_{ij}^- = f(x_1) + \sum_{ij \in \pi} |f_{ij}| = \Theta(f^\#, f)$ .

Finally, to each view point  $X$  corresponds a decomposition of the function  $f$  into a difference between a function representing the peaks and another representing the valleys:  $f = \Theta_X(e^+, f) - \Theta_X(e^-, f)$ .

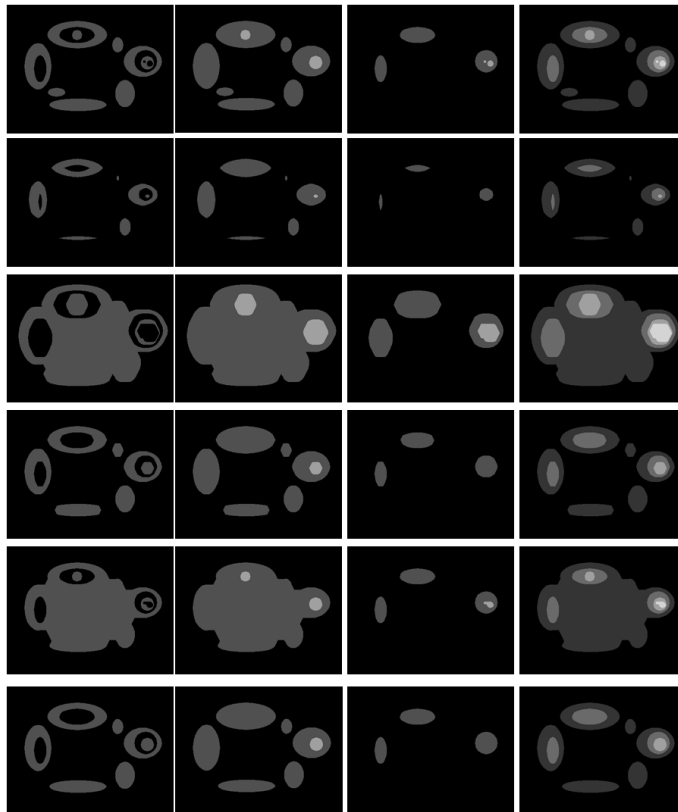
## 4 The lattice of floodings with a unique regional minimum

The schedules  $\widehat{\tau}_X(x) = \Theta_X(e, \tau)$  is the length of the shortest path between  $\Omega$  and the node  $x$ . The edge weights being non negative, there exists a non ascending path between any node and  $\Omega$  for the valuations  $\widehat{\tau}_X$ . Hence  $\Omega$  is the only regional minimum of  $\widehat{\tau}_X$ .

The following formulations are equivalent:

- $\Omega$  is the only regional minimum of  $\widehat{\tau}_X$
- the only regional minima of  $\tau$  belong to  $X$
- $\widehat{\tau}_X$  is invariant by the swamping which imposes  $\Omega$  as only regional minimum ; that is the flooding  $\Lambda^+(\widehat{\tau}_X, 0_\Omega^\infty)$ , where  $0_\Omega^\infty$  is equal to 0 at  $\Omega$  and  $\infty$  elsewhere.
- $\widehat{\tau}$  is invariant by the swamping which imposes the minima of  $X$  as only regional minima ; that is the flooding  $\Lambda^+(\widehat{\tau}, \tau_X^\infty)$
- considering the threshold at valuation  $\lambda$ , the subgraph spanning all nodes with a valuation  $\tau < \lambda$  (we call it the background at level  $\lambda$ ) has only one connected component, and this component contains  $\Omega$ .

We call  $F_X$  the lattice of all functions verifying the previous equivalent criteria.



**Fig. 2.** Decomposition and transformations on binary images

#### 4.1 Infimum and supremum in the lattice $F_X$

$F_X$  is a complete lattice, with the ordinary order relations for functions, 0 as minimal element,  $\infty$  as maximal element.

If  $h_1$  and  $h_2$  are two functions of  $F_X$ , their infimum in  $F_X$  is simply  $h_1 \wedge h_2$ . As a matter of fact, the backgrounds of  $h_1$  and  $h_2$  at level  $\lambda$  comprise each one connected component ; as they both contain  $\Omega$ , their union, which is the background of  $h_1 \wedge h_2$  also consists in one connected component containing  $\Omega$ .

Now,  $h_1 \vee h_2$  may have regional minima outside  $X$ , which have to be suppressed by swamping. For this reason the supremum  $h_1 \sqcup h_2$  is equal to  $\Lambda^+(h_1 \vee h_2; (h_1 \vee h_2)_X^\infty)$ .

**Remark:** These operators are not distributive one with another.

## 4.2 An adjunction in the lattice $F_X$

If for each node  $x$  we define a neighborhood  $V_x$ , the erosion of a function  $f$  with this variable structuring element is  $\varepsilon_V f(x) = \bigwedge_{y \in V_x} f_y$ . This operator is an erosion both in the ordinary lattice of functions as in the lattice  $F_X$ . Its adjunct dilation in the ordinary lattice is  $\delta_V f(y) = \bigvee_{x \in V_y^t} f_x$ , where  $V_y^t = \{x \mid y \in V_x\}$ . Being

an adjunction, they verify for each couple  $(f, g)$  of functions the relation:  $f < \varepsilon_V g \Leftrightarrow \delta_V f < g$ .

As  $\delta_V f(y) = \bigvee_{x \in V_y^t} f_x$  does generally not belong to  $F_X$ , the adjunct of  $\varepsilon_V g$  is

$\widetilde{\delta_V f}(y) = \bigsqcup_{x \in V_y^-} f_x = \Lambda^+(\delta_V f, (\delta_V f)_X^\infty)$ . To establish it we have to prove that

$\delta_V f < g \Leftrightarrow \Lambda^+(\delta_V f, (\delta_V f)_X^\infty) < g$ .

On one hand, floodings being extensive we have  $\Lambda^+(\delta_V f, (\delta_V f)_X) < g \Rightarrow \delta_V f < g$ .

On the other hand, if  $f$  and  $g$  belong to  $F_X$ , then  $\delta_V f < g \Rightarrow \Lambda^+(\delta_V f, (\delta_V f)_X) < \Lambda^+(g, g_X)$ . But  $\Lambda^+(g, g_X) = g$ , as  $g$  is invariant by the flooding  $\Lambda^+(g, g_X)$ .

Concatenating all equivalences, we get for any couple of functions in  $F_X$ :  $f < \varepsilon_V g \Leftrightarrow \delta_V f < g \Leftrightarrow \widetilde{\delta_V f} = \Lambda^+(\delta_V f, (\delta_V f)_X) < g$ , showing that  $(\varepsilon_V f, \widetilde{\delta_V f}) = (\bigwedge_{y \in V_x} f_y, \bigsqcup_{x \in V_y^-} f_x)$  is indeed an adjunction on  $F_X$ .

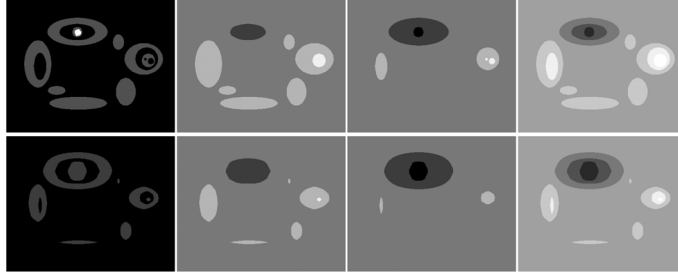
Classically  $\varepsilon_V \widetilde{\delta_V}$  is then a closing and  $\widetilde{\delta_V} \varepsilon_V$  an opening.

## 4.3 Razings in the lattice $F_X$

Razings clip peaks and do not create new holes. Thus the image of  $F_X$  by any type of razing or reconstruction opening is  $F_X$ .

## 5 Min, max and morphological operators through the decomposition in peaks and wells

Suming up, given a set  $X$ , serving as view-point, we are now able to decompose any image  $f$  into a cumulative image of its peaks  $f_X^+$  and a cumulative image of its wells  $f_X^-$ , such that  $f = f_X^+ - f_X^-$ . Both functions belong to  $F_X$ , a lattice where we defined a supremum, infimum and an adjunction. Furthermore this lattice is stable by any type of razing.



**Fig. 3.** Decomposition and transformations on binary images. The view point set is the white particle in the first image.

### 5.1 "Supremum" and "infimum" of two functions

Let  $(f_X^+, f_X^-)$  and  $(g_X^+, g_X^-)$  be the decomposition of two function into peaks and wells, given the view point  $X$ . For any composition law  $\Delta$  on  $F_X$ , we may define  $f \Delta g = f_X^+ \Delta g_X^+ - f_X^- \Delta g_X^-$ .

**Remark:** We do not necessarily have  $(f \Delta g)_X^+ = f_X^+ \Delta g_X^+$  and  $(f \Delta g)_X^- = f_X^- \Delta g_X^-$ . This means that the decomposition  $(f_X^+ \Delta g_X^+, f_X^- \Delta g_X^-)$  is not necessarily a minimal decomposition of  $f \Delta g$ .

Based on the operator  $\wedge$  and  $\sqcup$  in  $F_X$ , we obtain  $f \wedge g = f_X^+ \wedge g_X^+ - f_X^- \wedge g_X^-$  and  $f \vee g = f_X^+ \sqcup g_X^+ - f_X^- \sqcup g_X^-$ . Both operators are illustrated by figure 1, where the first row shows in position 1 and 2 the two initial binary sets  $X$  and  $Y$ , in position 3 their intersection  $f \wedge g$  and in position 4 their union  $f \vee g$ . Row 2 presents the image of peaks and row 3 the image of wells. These results are obtained by constructing the decomposition  $(f_X^+, f_X^-)$  and  $(g_X^+, g_X^-)$  which are then combined to obtain  $f \wedge g$  and  $f \vee g$ .

### 5.2 Operators operating on the peaks and wells

Given the decomposition  $(f_X^+, f_X^-)$  of a function  $f$  into peaks and wells, and given an increasing operator  $\psi$  from  $F_X$  into  $F_X$ , we define  $\hat{\psi}f = \psi(f_X^+) - \psi(f_X^-)$ . The decomposition is not necessarily minimal, but since  $f_X^- < f_X^+$  and  $\psi$  increasing,  $\hat{\psi}f$  will be positive if  $f$  is positive.

Figure 2 presents a number of operators on a binary image. The view point set  $X$  is the boundary of the image. In each row, the initial image presents a binary image, the second its peaks, the third its wells, and the last the total variation, that is the sum of peaks and wells. Each following row presents a particular transform  $\psi$ ; we construct  $\psi(f_X^+)$  and  $\psi(f_X^-)$  shown in columns 2 and 3. The subtraction  $\hat{\psi}f = \psi(f_X^+) - \psi(f_X^-)$  gives the resulting binary image in column 1 and the total variation  $\psi(f_X^+) + \psi(f_X^-)$  in column 4.

The transforms are presented row wise: 1) initial image; 2) erosion of size 11; 3) dilation of size 15; 4) opening of size 11; 5) closing of size 15; 6) reconstruction opening of size 11.





**Fig. 4.** Decomposition and transformation of grey tone images.

Figure 3 also presents an erosion of size 11 of the same binary image, but the decomposition is made with respect to a different view point set  $X$ , shown in white superimposition in the first image.  $X$  is contained in a white particle, itself contained in a hole contained in a particle. The disposition of the images is the same as for the preceding image. The grey tone scale is artificially enhanced so as to make all structures visible. Let us observe the decomposition. One has to think that the shortest paths start within  $X$  until they reach the boundary of the first particle, which appears as a downwards transition. They leave the first hole on an upwards transition, and again leave this particle on a downwards transition. Their paths stay in the background until the other particles are met. The transitions when they enter in the other particles are the same as for a path coming from the boundary of the image. This example shows the utmost

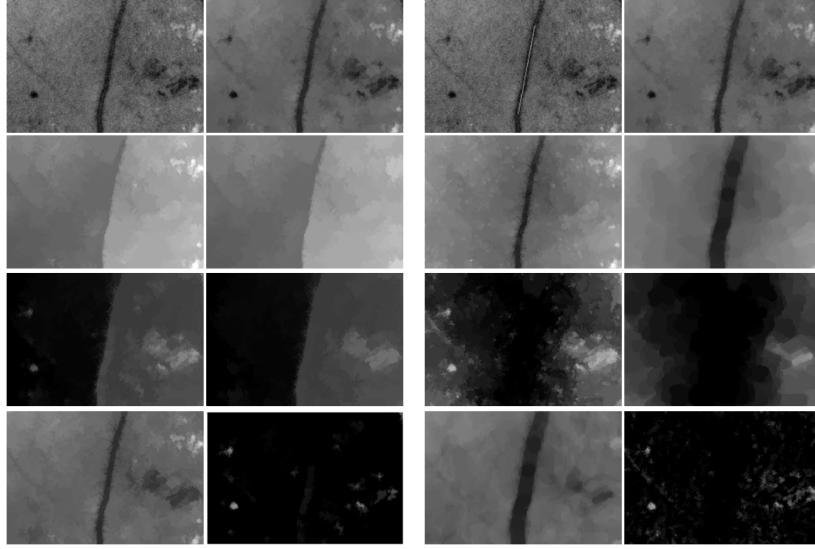
importance of the choice of  $X$ . Particles containing  $X$  are treated differently from particles which do not contain it. The second row shows the erosion. As the particles and holes containing  $X$  appear as holes both in the peak as in the wells image, their erosion enlarges them. On the contrary, particles and holes not containing  $X$  appear as peaks and are indeed eroded. The recomposition  $\widehat{\varepsilon}f = \varepsilon(f_X^+) - \varepsilon(f_X^-)$  is shown in image 2.1 and the total variation  $\varepsilon(f_X^+) + \varepsilon(f_X^-)$  in image 2.4.

Fig.4 presents the decomposition and transformation of grey tone images ; there are four groups with an identical disposition. In each group the upper row shows the positive variation  $\sigma^+$  and negative variation  $\sigma^-$ , whereas the second row shows the initial image and the total variation  $\sigma^\#$ . The first group presents the initial image and its decomposition. The second group the erosion of size 3, the third the dilation of size 3 and the last the opening of size 3. Since the first letter "P" of "Paulus" intersects the boundary of the image, which plays the role of view point set  $X$ , it is not transformed as the other letters.

The last example in fig. 5 is taken from an image of the retina. The aim is to detect the small dots and the bleedings without detecting the vessel, although they have the same width. The images are organized as two matrices with the same structure. For the left block the view point set is the left boundary of the image, and for the right block the white segment in superimposition with the vessel in image 1.1 Both blocks are organized as follows. Fig.1.1 the initial image, 1.2 its filtered version which is decomposed. The peaks are in fig 2.1, the wells in fig 3.1. The tranformation applied on peaks and wells in the left block is a reconstruction opening and in the right block an erosion. The final result of the transform is in fig. 4.1. Subtracting image 1.2 from the maximum between 1.2 and 4.1 yields a residue image where the anevrysms are clearly highlighted and the vessel disappeared. Chosing the left boundary of the image as view point  $X$  for the left block produces a decomposition where the vessel is split: its right boundary appear in the peak image and its left boundary in the wells image, whereas the microanevrysms appear a dots in the wells image. For this reason, the reconstruction opening leaves the vessel untouched and suppresses the small dots. In the right block, the image is seen from a point view image included in the vessel. For this reason the boundaries of the vessel appear as upwards transitions and the boundaries of the microanevrysms as downward transitions. Like that they completely diverge in the treatment: the vessels is eroded in the peaks image and the dots disappear by eroding the wells image. The recomposed image shows a blown up vessel and suppressed anevrysms.

## 6 Conclusion

The decomposition we proposed depends on the edge and node weights and upon the point of view, the set  $X$ . Changing  $X$  produces a different decomposition of the same scene, with a different distribution of contrast, between the same objects, which may facilitate the segmentation. Changing the edge weights, for instance by keeping only the variations along each edge which are higher than



**Fig. 5.** Analysis of the retina in order to detect the micro-anevrysms and discard the vessel

a threshold will put the focus with contrasted contours and ignore the objects with a more fuzzy contour. The distribution of weights on the edges may also be directional, favoring some directions and discarding others. This gives the possibility to analyze anisotropies. The shortest distances which are at the basis of the decomposition may be rendered geodesic. In conclusion, the framework set in this paper should be useful for many tasks in image processing: filtering, segmentation and pattern recognition in general.

## References

1. Coloma Ballester, Vicent Caselles, and P. Monasse. The tree of shapes of an image. *ESAIM: COCV*, 9:1–18, 2003.
2. Henk J. A. M. Heijmans. Connected morphological operators for binary images. *Computer Vision and Image Understanding*, 73(1):99 – 120, 1999.
3. Renato Keshet. Adjacency lattices and shape-tree semilattices. *Image and Vision Computing*, 25(4):436 – 446, 2007. International Symposium on Mathematical Morphology 2005.
4. Ch. Ronse. *Bounded variation in posets, with applications in morphological image processing*, volume 86, pages 249–283. Acta Universitatis Upsaliensis, 2009. Proceedings of the Kiselmanfest.
5. J. Serra. *Mathematical Morphology. Volume I*. London: Academic Press, 1982.